

61A LECTURE 25 – DECLARATIVE PROGRAMMING

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Midterm grades up

- Class did well as a whole!
- `gllookup -s test2`
- Regrades: Talk to your TA in person.

Albert keeps all of his top secret information in a binary tree. This prevents the layperson from reading his data. However, well trained computer scientists (such as you) can still access his information.

As a further layer of protection, Albert turns some of the nodes in his trees into **Eert** nodes. **Eert** nodes, which have **Tree** as their base class, are like normal **Tree** nodes, except they swap their left and right branches. (Albert settles for nothing less than the most advanced encryption techniques known to man.)

- (a) (3 pt) Complete the `__init__` method for the **Eert** class on the next page. Make sure to use inheritance as much as possible. The **Eert** class should work as follows:

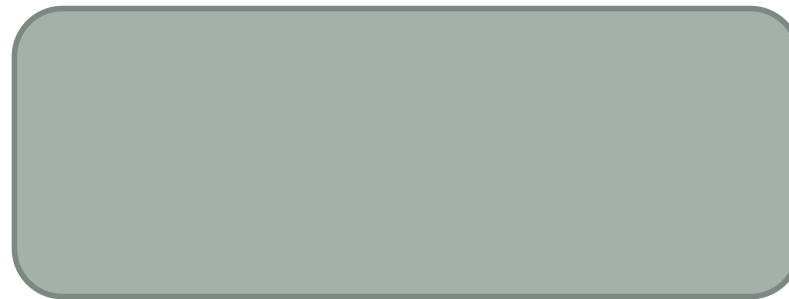
```
>>> e = Eert("61A account info",
...         Tree("Username: cs61a-te"),
...         Tree("Password: imsocool"))
>>> e.entry # unchanged
"61A account info"
>>> e.left.entry # swapped with right
"Password: imsocool"
>>> e.right.entry # swapped with left
"Username: cs61a-te"
```

- (b) (5 pt) Complete the definitions of the `decrypt` methods for both the **Tree** and **Eert** classes on the next page. When the `decrypt` method is invoked on a binary tree containing **Tree** and **Eert** nodes, it returns a copy of the binary tree, but with all **Eert** nodes replaced with **Tree** nodes. During this replacement, you should also swap the left and the right back to their proper positions! Here is a graphical representation of the process:

```
class Tree(object):

    def __init__(self, entry, left=None, right=None):
        self.entry = entry
        self.left = left
        self.right = right

    def decrypt(self):
        """ PART B """
```



```
class Eert(Tree):

    def __init__(self, entry, left=None, right=None):
        """ PART A """
```



```
    def decrypt(self):
        """ PART B """
```



Announcements

- Proj4 has been out
 - Due in 7 days – Start if you haven't!
 - Recursive art contest deadline one day **before** project is due
 - Future “homework” assignment will be to vote on your favorite submissions
- Final exam next Thursday
 - 7-10pm in 1 Pimentel
 - Any conflicts – notify us immediately
- Final exam review session this weekend
 - See Piazza Poll to vote on your time
 - Potential extra credit – more information later in the week

Laziness

Recall our previous sequence interface:

- A sequence has a finite, known length
- A sequence allows element selection for any element

In the cases we've seen so far, satisfying the sequence interface requires storing the entire sequence in a computer's memory

Problems?


- Infinite sequences - primes, positive integers
- Really large sequences - all Twitter tweets, votes in a presidential election

Streams

A stream is a recursive list with an *explicit* first element and a *lazily computed* rest-of-the-list

```
class Stream(Rlist):
    """A lazily computed recursive list."""
    def __init__(self, first,
                  compute_rest=lambda: Stream.empty):
        assert callable(compute_rest)
        self.first = first
        self._compute_rest = compute_rest
        self._rest = None

    @property
    def rest(self):
        """Return the rest of the stream, computing it if
        necessary."""
        if self._compute_rest is not None:
            self._rest = self._compute_rest()
            self._compute_rest = None
        return self._rest
```



Integer Streams

An integer stream is a stream of consecutive integers

An integer stream starting at k consists of k and a function that returns the integer stream starting at $k+1$

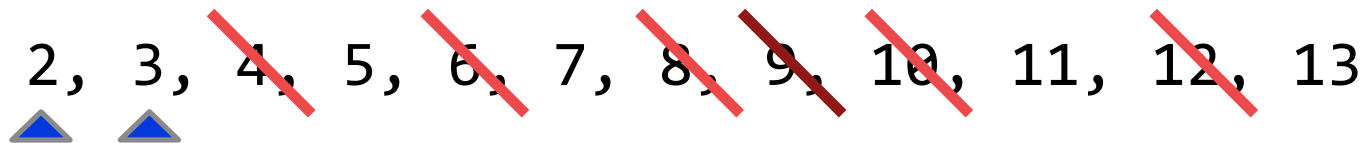
```
def integer_stream(first=1):  
    """Return a stream of consecutive integers, starting  
    with first.  
  
    >>> s = integer_stream(3)  
    >>> s.first  
    3  
    >>> s.rest.first  
    4  
    """  
    def compute_rest():  
        return integer_stream(first+1)  
    return Stream(first, compute_rest)
```


A Stream of Primes

The stream of integers not divisible by any $k \leq n$ is:

- The stream of integers not divisible by any $k < n$,
- Filtered to remove any element divisible by n
- This recurrence is called the *Sieve of Eratosthenes*

2, 3, ~~4~~, 5, ~~6~~, 7, ~~8~~, ~~9~~, ~~10~~, 11, ~~12~~, 13



```
def primes(istream):  
    """Return a stream of primes, given a stream of  
    consecutive integers."""  
    def compute_rest():  
        not_divisible = lambda x: x % istream.first != 0  
        return primes(filter_stream(not_divisible,  
                                    istream.rest))  
    return Stream(istream.first, compute_rest)
```

Try it

- Write a function `add_streams` that takes two streams and returns a new stream formed by summing corresponding elements in the argument streams. Stop when either of the streams ends.
- Bonus: see if you can use `add_streams` to define the Fibonacci stream!

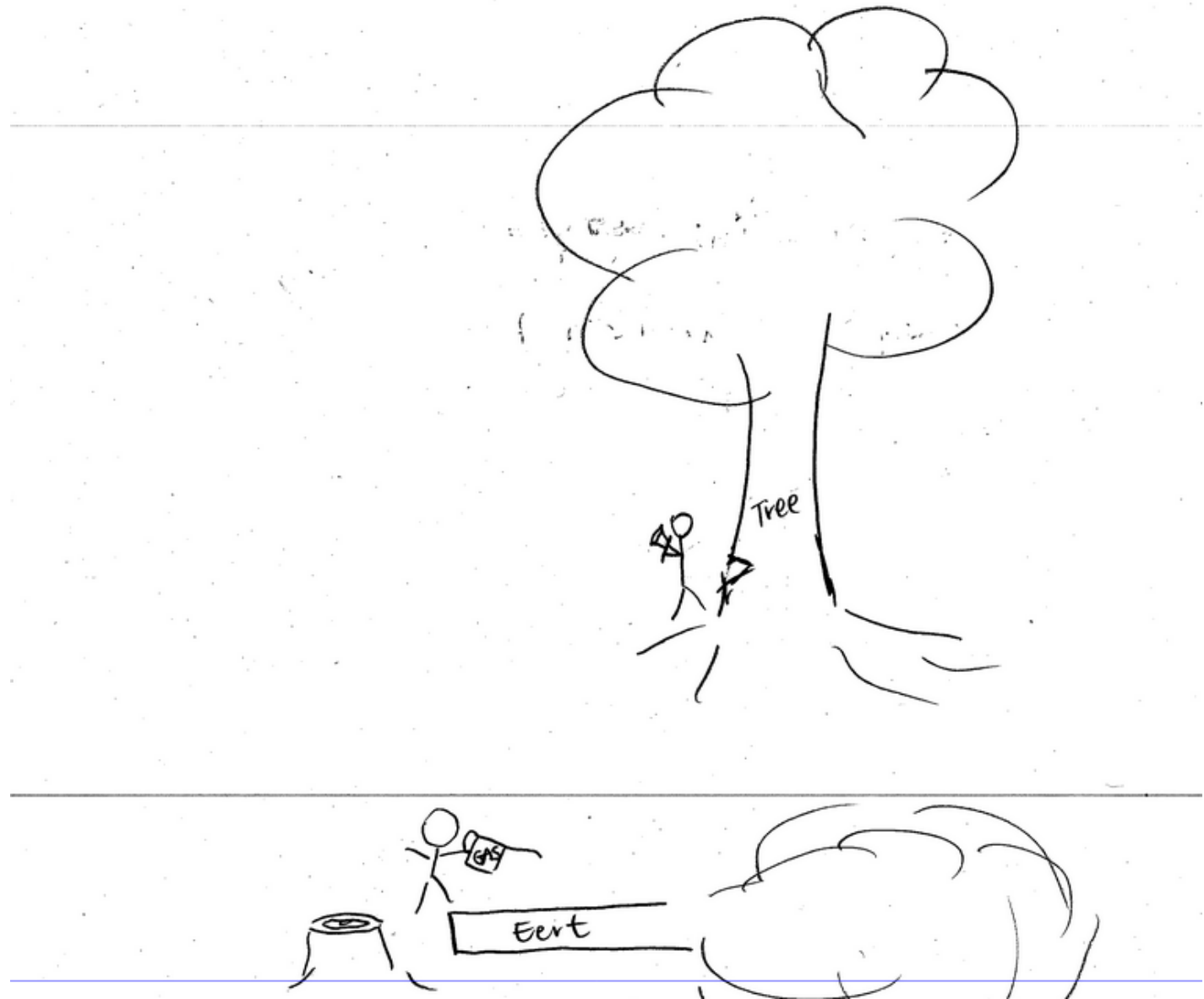
Answers

```
def add_streams(s1, s2):  
    if s1 is Stream.empty or s2 is Stream.empty:  
        return Stream.empty  
    return Stream(s1.first + s2.first,  
                  lambda: add_streams(s1.rest, s2.rest))  
  
fibs = Stream(0,  
              lambda: Stream(1,  
                              lambda: add_streams(fibs, fibs.rest)))
```

Short Break

8. (0 points) Express yourself (v2)

Express your feelings in the space below through your choice of creative medium, such as poetry or illustration.



Last “super-big” topic in course

- Lot to cover this lecture...
- We will continue this topic tomorrow as well
- Need to finish questions 1-4 on your Scheme project for Logic programming to work
- Bring your version of scheme to lab tomorrow!

Databases

A database is a collection of records (tuples) and an interface for adding, editing, and retrieving records

The Structured Query Language (SQL) is perhaps the most widely used programming language on Earth

```
SELECT * FROM toy_info WHERE color='yellow';
```

| toy_id | toy | color | cost | weight |
|--------|-------------|--------|------|--------|
| 2 | whiffleball | yellow | 2.20 | 0.40 |
| 5 | frisbee | yellow | 1.50 | 0.20 |
| 10 | yoyo | yellow | 1.50 | 0.20 |

SQL is an example of a declarative programming language.

It separates *what* to compute from *how* it is computed

The language interpreter is free to compute the result in any way it deems appropriate

http://www.headfirstlabs.com/sql_hands_on/

Declarative Programming

The main characteristics of declarative languages:

- A "program" is a description of the desired solution
- The interpreter figures out how to generate such a solution

By contrast, in procedural languages such as Python & Scheme:

- A "program" is a description of procedures
- The interpreter carries out execution/evaluation rules

Building a universal problem solver is a difficult task

Declarative programming languages compromise by solving only a subset of all problems

They typically trade off data scale for problem complexity

The Logic Language

The *Logic* language is invented for this course

- Based on the Scheme project & ideas from Prolog
- Expressions are facts or queries, which contain relations
- Expressions and relations are both Scheme lists
- For example, (**likes Albert dogs**) is a relation
- Implementation fits on a single sheet of paper

Today's theme:



<http://awhimsicalbohemian.typepad.com/.a/6a00e5538b84f3883301538dfa8f19970b-800wi>

Simple Facts

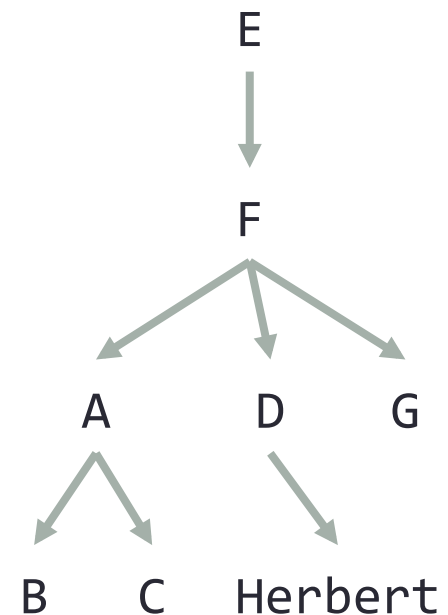
A simple fact expression in the *Logic* language declares a relation to be true

Let's say I want to track my many dogs' ancestry

Language Syntax:

- A relation is a Scheme list
- A fact expression is a Scheme list containing **fact** followed by one or more relations

```
logic> (fact (parent delano herbert))  
logic> (fact (parent abraham barack))  
logic> (fact (parent abraham clinton))  
logic> (fact (parent fillmore abraham))  
logic> (fact (parent fillmore delano))  
logic> (fact (parent fillmore grover))  
logic> (fact (parent eisenhower fillmore))
```



Relations are Not Procedure Calls

In *Logic*, a relation is not a call expression

- In Scheme, we write **(abs -3)** to call **abs** on **-3**
- In *Logic*, **(abs -3 3)** asserts that the **abs** of **-3** is **3**

For example, if we wanted to assert that **1 + 2 = 3**:

(add 1 2 3)

Why declare knowledge in this way? It will allow us to solve problems in two directions:

(add 1 2 _)

(add _ 2 3)

(add 1 _ 3)

(___ 1 2 3)

Queries

A query contains one or more relations. The *Logic* interpreter returns whether (and how) they are all simultaneously satisfied

Queries may contain variables: symbols starting with **?**

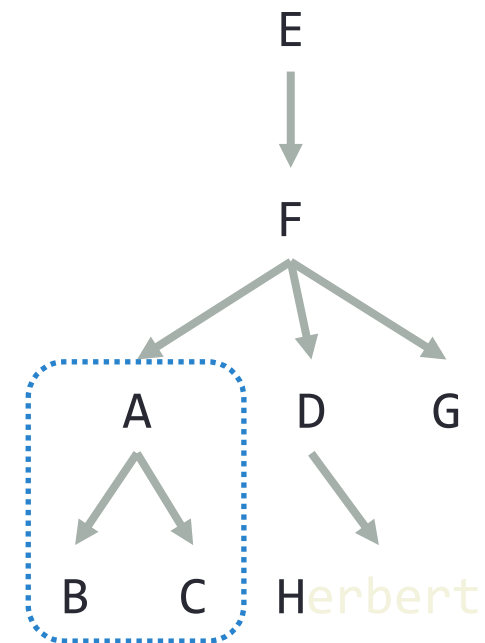
```
logic> (fact (parent delano herbert))
logic> (fact (parent abraham barack))
logic> (fact (parent abraham clinton))
logic> (fact (parent fillmore abraham))
logic> (fact (parent fillmore delano))
logic> (fact (parent fillmore grover))
logic> (fact (parent eisenhower fillmore))
```

```
logic> (query (parent abraham ?child))
```

Success!

child: barack

child: clinton



Queries

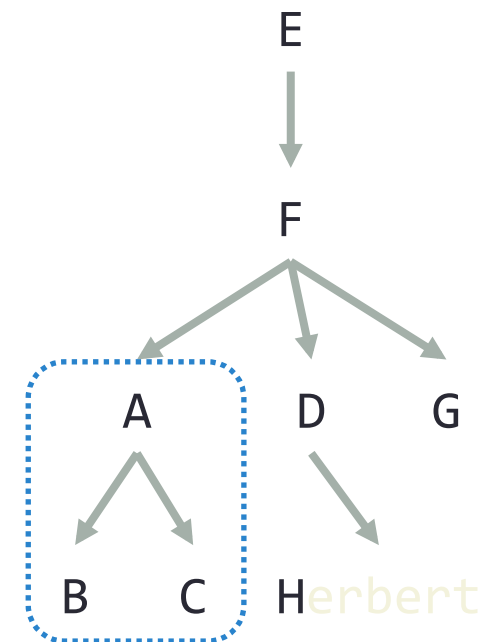
A query contains one or more relations. The *Logic* interpreter returns whether (and how) they are all simultaneously satisfied

Queries may contain variables: symbols starting with ?

```
logic> (fact (parent delano herbert))
logic> (fact (parent abraham barack))
logic> (fact (parent abraham clinton))
logic> (fact (parent fillmore abraham))
logic> (fact (parent fillmore delano))
logic> (fact (parent fillmore grover))
logic> (fact (parent eisenhower fillmore))

logic> (query (parent ?who barack)
              (parent ?who clinton))
```

Success!
who: abraham



Compound Facts

A fact can include multiple relations and variables as well

(fact <conclusion> <hypothesis₀> <hypothesis₁> ... <hypothesis_N>)

Means <conclusion> is true if all <hypothesis_k> are true

```
logic> (fact (child ?c ?p) (parent ?p ?c))
```

```
logic> (query (child herbert delano))
```

Success!

```
logic> (query (child eisenhower clinton))
```

Failure.

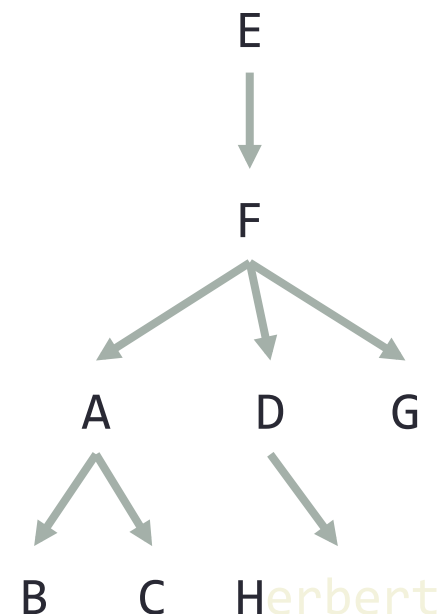
```
logic> (query (child ?child fillmore))
```

Success!

child: abraham

child: delano

child: grover



Recursive Facts

A fact is recursive if the same relation is mentioned in a hypothesis and the conclusion

```
logic> (fact (ancestor ?a ?y) (parent ?a ?y))
```

```
logic> (fact (ancestor ?a ?y) (parent ?a ?z) (ancestor ?z ?y))
```

```
logic> (query (ancestor ?a herbert))
```

Success!

a: delano

a: fillmore

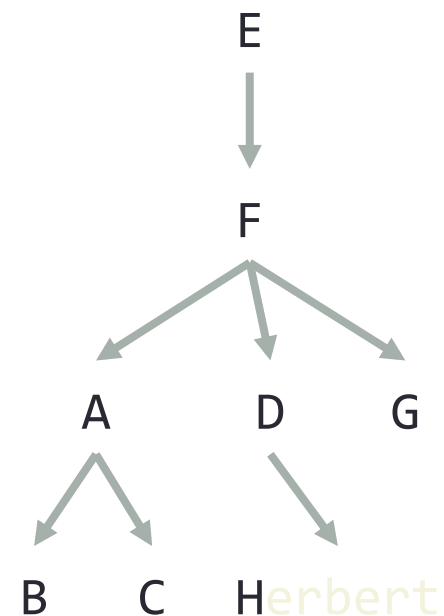
a: eisenhower

```
logic> (query (ancestor ?a barack)
              (ancestor ?a herbert))
```

Success!

a: fillmore

a: eisenhower



Searching to Satisfy Queries

The *Logic* interpreter performs a search in the space of relations for each query to find a satisfying assignment

```
logic> (query (ancestor ?a herbert))
```

Success!

a: delano

a: fillmore



a: eisenhower

```
logic> (fact (parent delano herbert))
```

```
logic> (fact (parent fillmore delano))
```

```
logic> (fact (ancestor ?a ?y) (parent ?a ?y))
```

```
logic> (fact (ancestor ?a ?y) (parent ?a ?z) (ancestor ?z ?y))
```

(parent delano herbert) ; (1), a simple fact

(ancestor delano herbert) ; (2), from (1) and the 1st ancestor fact

(parent fillmore delano) ; (3), a simple fact

(ancestor fillmore herbert) ; (4), from (2), (3), & the 2nd ancestor fact

Hierarchical Facts

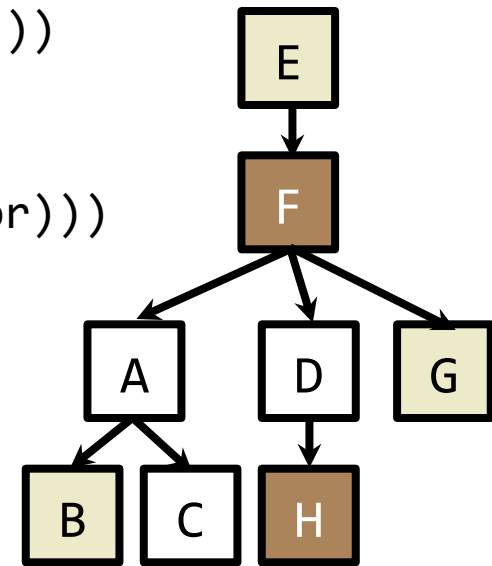
Relations can contain relations in addition to atoms

```
logic> (fact (dog (name abraham) (color white)))  
logic> (fact (dog (name barack) (color tan)))  
logic> (fact (dog (name clinton) (color white)))  
logic> (fact (dog (name delano) (color white)))  
logic> (fact (dog (name eisenhower) (color tan)))  
logic> (fact (dog (name fillmore) (color brown)))  
logic> (fact (dog (name grover) (color tan)))  
logic> (fact (dog (name herbert) (color brown)))
```

Variables can refer to atoms or relations

```
logic> (query (dog (name clinton) (color ?color)))  
Success!  
color: white
```

```
logic> (query (dog (name clinton) ?info))  
Success!  
info: (color white)
```



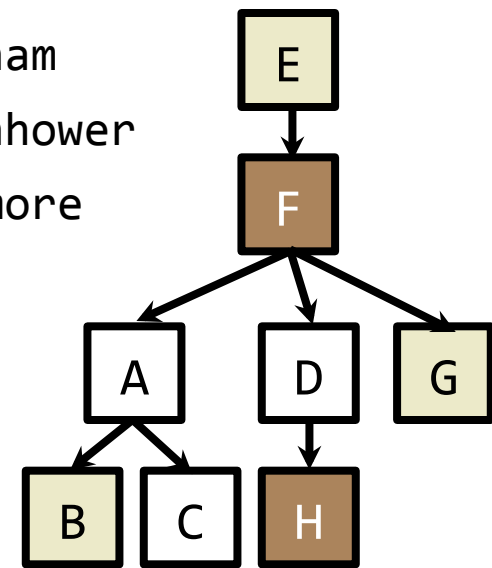
Example: Combining Multiple Data Sources

Which dogs have an ancestor of the same color?

```
logic> (query (dog (name ?name) (color ?color))  
            (ancestor ?ancestor ?name)  
            (dog (name ?ancestor) (color ?color)))
```

Success!

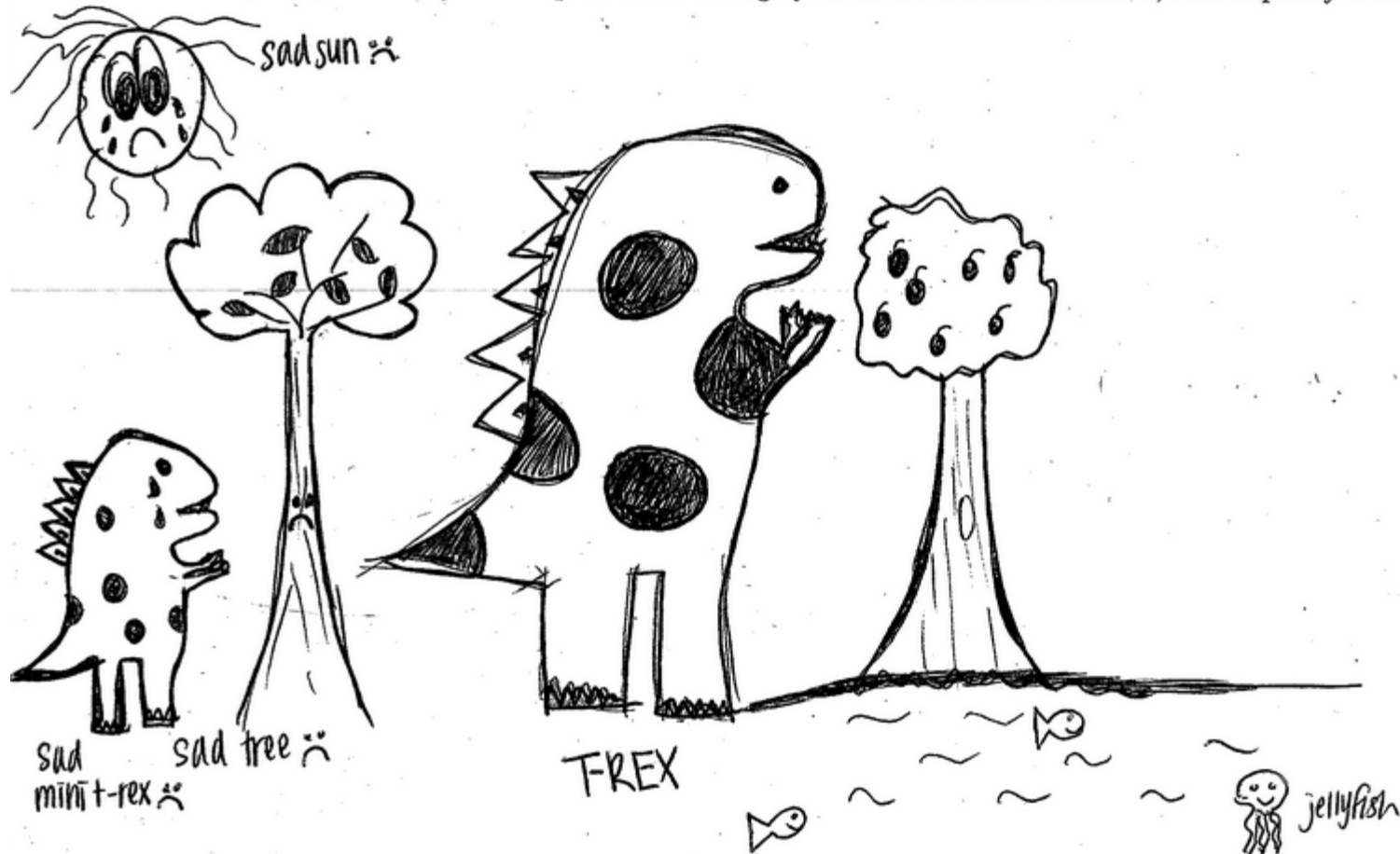
| | | |
|---------------|--------------|----------------------|
| name: barack | color: tan | ancestor: eisenhower |
| name: clinton | color: white | ancestor: abraham |
| name: grover | color: tan | ancestor: eisenhower |
| name: herbert | color: brown | ancestor: fillmore |



Break

8. (0 points) Express yourself (v2)

Express your feelings in the space below through your choice of creative medium, such as poetry or illustration.



The berrytree is always bluer in somebody else's lake.
T-rex dreams about going out there, but that is a big mistake.

Example: Appending Lists

Two lists append to form a third list if:

- The first list is empty and the second and third are the same

`() (a b c) (a b c)`

- Both of the following hold:
 - List 1 and 3 have the same first element
 - The rest of list 1 and all of list 2 append to form the rest of list 3

`(a b c) (d e f) (a b c d e f)`

```
logic> (fact (append-to-form () ?x ?x))
```

```
logic> (fact (append-to-form (?a . ?r) ?y (?a . ?z))  
          (append-to-form ?r ?y ?z))
```

Logic Example: Anagrams

A permutation (i.e., anagram) of a list is:

- The empty list for an empty list
- The first element of the list inserted into an anagram of the rest of the list

Element

List

List with element

```
(fact (insert ?a ?r ((?a . ?r))))
```

```
(fact (insert ?a (?b . ?r) (?b . ?s))  
      (insert ?a ?r ?s))
```

```
(fact (anagram () ()))
```

```
(fact (anagram (?a . ?r) ?b)  
      (insert ?a ?s ?b)  
      (anagram ?r ?s))
```

a | r t

r t

a r t

r a t

r t a

t r

a t r

t a r

t r a

You try it out!

- Write facts to make double-elements work

```
logic> (query (double-elements (3 4) ?result))
```

Success!

```
result: (3 3 4 4)
```

```
logic> (query (double-elements ?start (4 4 5 5)))
```

Success!

```
start: (4 5)
```